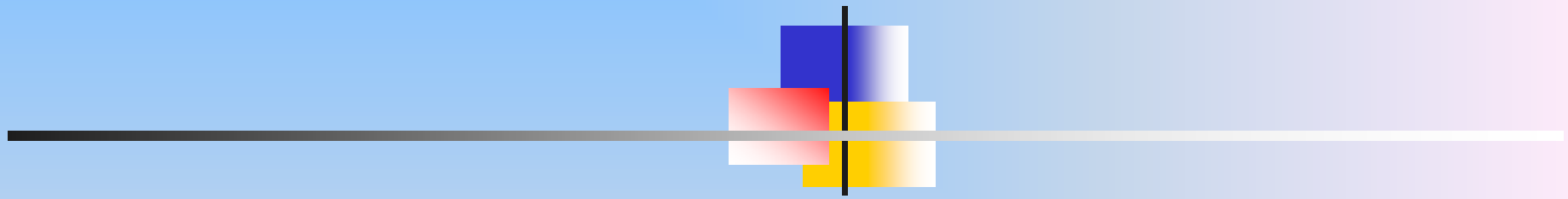
CSE302: Compiler Design



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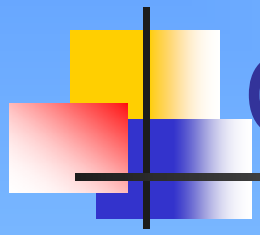
April 19, 2007



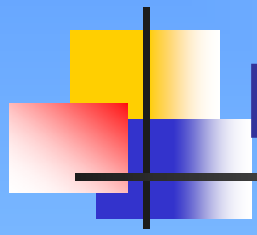
Outline

- Recap
  - Intermediate code generation
- Things related to project part II
- Summary

Translation of Boolean Expressions (1)

- S→if(B) S1 {B.true=newlabel(); B.false=S.next;} {S.code=B.code || label(B.true) || S1.code
  - if B goto L1
  - goto L2
  - L1: S1
  - L2:
- B → B1 || B2
- S → if(B1||B2) S1 {B.true=newlabel(); B.false=S.next; B1.true=B.true; B1.false=newlabel(); B2.true=B.true; B2.false=B.false;} {B.code=B1.code || label(B1.false) || B2.code} {S.code=B.code || label(B.true) || S1.code
  - if B1 goto L1
  - goto L2
  - L2: if B2 goto L1
  - goto L3
  - L1: S1
  - L3:

Translation of Boolean Expressions (2)



- B → B1 && B2
- S → if({ B.true=newlabel(); B.false=S.next;
  B1.true=newlabel(); B1.false=B.false;
  B2.true=B.true; B2.false=B.false;}
  B1&&B2) {B.code=B1.code || label(B1.true) || B2.code}
  S1 {S.code=B.code || label(B.true) || S1.code
  - if B1 goto L2
  - goto L3
  - L2: if B2 goto L1
  - goto L3
  - L1: S1
  - L3:

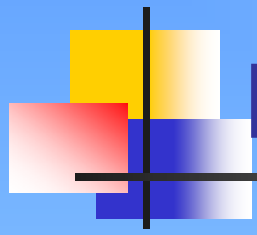
Translation of Boolean Expressions (3)

- B → !B1
- S → if({B.true=newlabel(); B.false=S.next;
  B1.true=B.false; B1.false=B.true;}
  !B1) {B.code=B1.code}
  S1 {S.code=B.code || label(B.true) || S1.code}
  - if B1 goto L2
  - goto L1
  - L1: S1
  - L2:

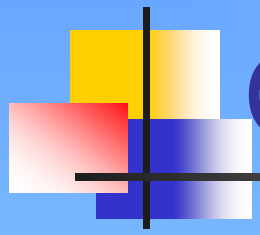
Translation of Boolean Expressions (4)

- B → E1 rel E2
- S → if(E1 rel E2) {B.true=newlabel(); B.false=S.next;}
  {B.code=E1.code || E2.code ||
  gen('if E1.addr rel.op E2.addr 'goto' B.true ||
  gen('goto' B.false);
  }
  S1 {S.code=B.code || label(B.true) || S1.code}
  - E1.code
  - E2.code
  - if E1 rel.op E2 goto L1
  - goto L2
  - L1: S1
  - L2:

Translation of Boolean Expressions (5)



- B $\rightarrow$ true {B.code=gen(‘goto’ B.true);}
- B $\rightarrow$ false {B.code=gen(‘goto’ B.false);}



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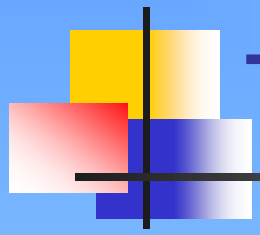
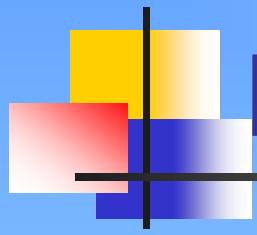
Support for Arbitrary Value Types



- YYSTYPE may be a union type

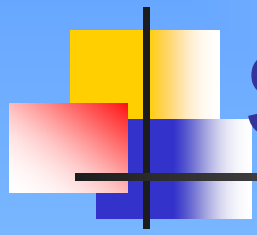
```
%union  {
    body of union ...
}
```

  - This declares the Yacc value stack, and the external variables yylval and yyval, to have type equal to this union.
  - -d: the union declaration is copied onto the y.tab.h
- The union member names are associated with terminals and nonterminals
  - %token <name> _TOKEN
  - %type <name> _nonterminal



LR parsing with L-attributes

- Section 5.5.4
- Yacc implementation

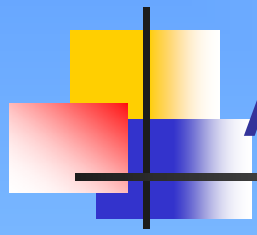


SDT for While

- SDD of **while** to generate 3-address code
  - $S \rightarrow$ while (C) S
    - while(i<a) i=i*2;
      j=j*i;
  - label L5: if i<a goto L6
    goto L7
    label L6: i=i*2
    goto L5
    label L7: j=j*i

- $S \rightarrow$ while (C) S L1=new(); L2=new();
  S1.next=L1;
  C.true=L2; C.false=S.next;

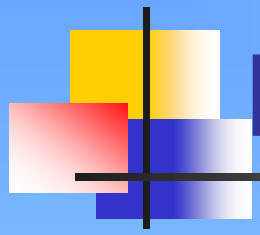
S.code= **label** || L1 || C.code || **label** || L2 || S1.code



Attribute Computation

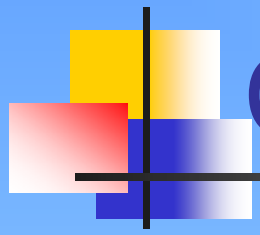
- For synthesized attribute computations
  - End of the production body
- For computing inherited attributes of a nonterminal
  - Immediately before the nonterminal occurrence

$S \rightarrow$ while ({ L1=new(); L2=new();
C.true=L2; C.false=S.next;
}
C) { S1.next=L1; }
S1 { S.code= **label** || L1 || C.code || **label** || L2 || S1.code; }



Markers

- $M \rightarrow \varepsilon$
  - Hold the inherited attributes



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