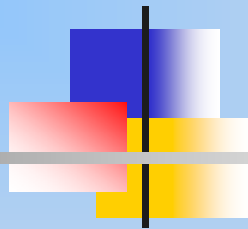
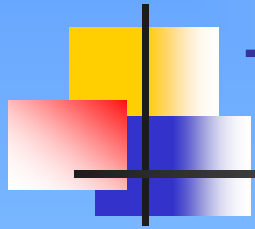


CSE302: Compiler Design



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January 23, 2007



Today's Outline

- Recap
 - Introduction (Section 2.1)
 - Syntax definition (Section 2.2)
 - Parsing (Section 2.4)
- A simple syntax-directed translator (Chapter 2)
 - Parsing (Section 2.4.5)
 - Syntax directed translation (Section 2.3)
 - A translator for simple expressions (Section 2.5)
- Summary and homework

BNF Grammar and Parse Trees

BNF Functionality

- Describing Lists
- Grammar & Derivation
- Parse Trees
- Avoiding Ambiguity

$\langle \text{program} \rangle \rightarrow \langle \text{stmts} \rangle$

$\langle \text{stmts} \rangle \rightarrow \langle \text{stmt} \rangle \mid \langle \text{stmt} \rangle ; \langle \text{stmts} \rangle$

$\langle \text{stmt} \rangle \rightarrow \langle \text{var} \rangle = \langle \text{expr} \rangle$

$\langle \text{var} \rangle \rightarrow a \mid b \mid c \mid d$

$\langle \text{expr} \rangle \rightarrow \langle \text{term} \rangle + \langle \text{term} \rangle \mid \langle \text{term} \rangle - \langle \text{term} \rangle$

$\langle \text{term} \rangle \rightarrow \langle \text{var} \rangle \mid \text{const}$

$\langle \text{program} \rangle$

$\Rightarrow \langle \text{stmts} \rangle$

$\Rightarrow \langle \text{stmt} \rangle$

$\Rightarrow \langle \text{var} \rangle = \langle \text{expr} \rangle$

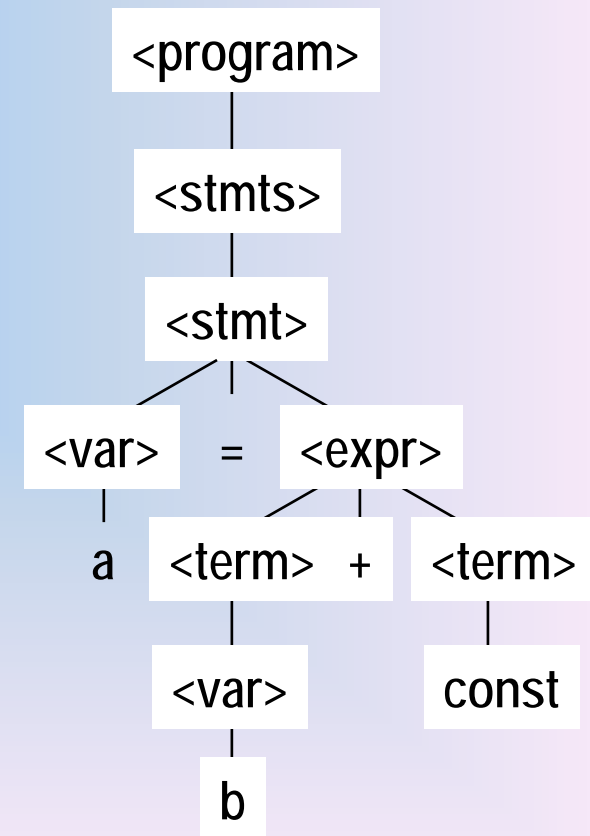
$\Rightarrow a = \langle \text{expr} \rangle$

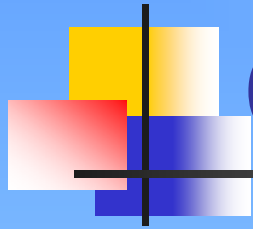
$\Rightarrow a = \langle \text{term} \rangle + \langle \text{term} \rangle$

$\Rightarrow a = \langle \text{var} \rangle + \langle \text{term} \rangle$

$\Rightarrow a = b + \langle \text{term} \rangle$

$\Rightarrow a = b + \text{const}$



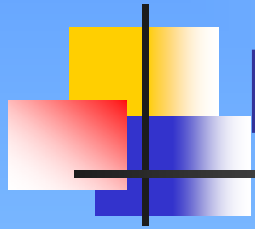


Grammar Ambiguity

BNF Functionality

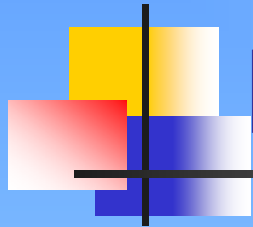
- Describing Lists
- Grammar & Derivation
- Parse Trees
- Avoiding Ambiguity

- A grammar is **ambiguous** iff it generates a sentential form that has two or more distinct parse trees
- Use BNF to specify operator precedence and associativity



Language Design

- Design a BNF grammar for a language that could express a one-digit number, an addition of two one-digit numbers, or a subtraction of two one-digit numbers
 - $\langle \text{expr} \rangle \rightarrow \langle \text{term} \rangle + \langle \text{term} \rangle \mid \langle \text{term} \rangle - \langle \text{term} \rangle \mid \langle \text{term} \rangle$
 - $\langle \text{term} \rangle \rightarrow 0 \mid 1 \mid 2 \mid \dots \mid 9$



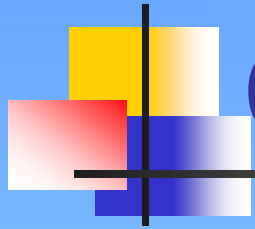
Language Implementation

- A recursive-descent parser
 - Language implementation **directly following the BNF grammar**
 - $\langle \text{expr} \rangle \rightarrow \langle \text{term} \rangle + \langle \text{term} \rangle \mid \langle \text{term} \rangle - \langle \text{term} \rangle \mid \langle \text{term} \rangle$
 - $\langle \text{term} \rangle \rightarrow 0 \mid 1 \mid 2 \mid \dots \mid 9$

- Pseudo code

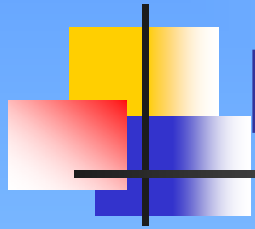
```
void expr() {  
    term();  
    if( token==plus_op  
        or token==minus_op) {  
        match(token);  
        term();  
    }  
    else error();  
}
```

```
void term() {  
    match(int_literal);  
}  
  
void match(expectedToken) {  
    if(token==expectedToken)  
        getNextToken();  
    else error();  
}
```



Outline

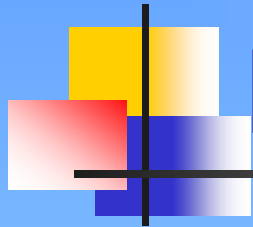
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Remove Left Recursion

- What are the languages defined by the following two BNF grammars?

$$A \rightarrow A \alpha \mid \beta$$
$$A \rightarrow \beta R$$
$$R \rightarrow \alpha R \mid \varepsilon$$



Remove Left Recursion

- BNF: $\langle \text{expr} \rangle \rightarrow \langle \text{expr} \rangle + \langle \text{term} \rangle$

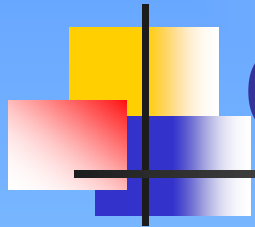
- Left-recursion to right-recursion

$$\begin{array}{c|c} A \rightarrow A \alpha \mid \beta & A \rightarrow \beta R \\ & R \rightarrow \alpha R \mid \varepsilon \end{array}$$

- $\langle \text{expr} \rangle \rightarrow \langle \text{term} \rangle \text{ rest}$

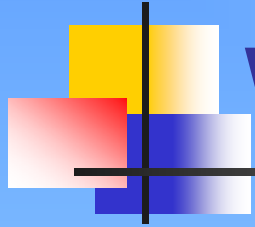
- $\text{rest} \rightarrow + \langle \text{term} \rangle \text{ rest} \mid - \langle \text{term} \rangle \text{ rest} \mid \varepsilon$

- EBNF: $\langle \text{expr} \rangle \rightarrow \langle \text{term} \rangle \{ + \langle \text{term} \rangle \}$



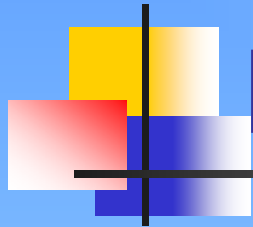
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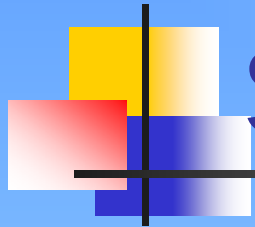
What Can Be Done So Far?

- Define language **syntax** using BNF grammar
- Parsing to detect **syntax** errors
 - Syntax analysis
- How about translation?
 - Syntax-directed translation
 - Attaching rules or program fragments to productions in a grammar
 - An example of translating infix notation to postfix notation



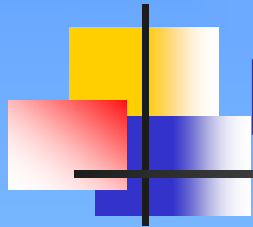
Postfix Notation

- Inductive definition
 - If E is a variable or constant, then the postfix notation for E is E itself
 - If E is an expression of the form $E1 \text{ op } E2$, then the postfix notation for E is $E1' E2' \text{ op}$
 - If E is of the form $(E1)$, then the postfix notation is $E1'$
 - Examples
 - The postfix notation $(9-5)+2$ is $95-2+$



Syntax-Directed Definition

- For a BNF grammar
 - Associate each grammar symbol (terminals and non-terminals) with a set of attribute
 - Type information for type checking/conversion
 - Notation representation for notation translation
 - Attach a **semantic rule** or **program fragment** to each production in a grammar
 - Computing the values of the attributes associated with the symbols in the production
- The BNF grammar becomes an **attribute grammar**



Definition of Attribute Grammar

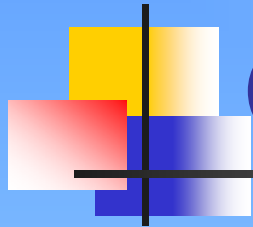
- An **attribute grammar** is a BNF grammar with additions:
 - For each grammar symbol x : a set $A(x)$ of **attribute** values
 - Each production in the grammar has a set of **semantic rules** that define or compute certain attributes of the nonterminals in the production
 - Each production in the grammar has a (possibly empty) set of **predicates** to check for attribute consistency
- A sentence derivation
 - Based on BNF
 - A parse tree
- A sentence derivation
 - Based on an attribute grammar
 - A fully attributed parse tree



A Type Checking Example Using Syntax-Directed Definition

- A BNF grammar
 - $\langle \text{assign} \rangle \rightarrow \langle \text{var} \rangle = \langle \text{expr} \rangle$
 - $\langle \text{expr} \rangle \rightarrow \langle \text{var} \rangle + \langle \text{var} \rangle$
 - $\langle \text{var} \rangle \rightarrow A \mid B \mid C$
- An attribute grammar
 1. Syntax production: $\langle \text{assign} \rangle \rightarrow \langle \text{var} \rangle = \langle \text{expr} \rangle$
 - Semantic rule: $\langle \text{expr} \rangle.\text{expected_type} \leftarrow \langle \text{var} \rangle.\text{actual_type}$
 2. Syntax production: $\langle \text{expr} \rangle \rightarrow \langle \text{var} \rangle + \langle \text{var} \rangle$
 - Semantic rule: $\langle \text{expr} \rangle.\text{actual_type} \leftarrow$

$\text{if}(\langle \text{var} \rangle[2].\text{actual_type} == \text{int}) \text{ and}$
 $(\langle \text{var} \rangle[3].\text{actual_type} == \text{int})$
 then int
 else real
 endif
 - Predicate: $\langle \text{expr} \rangle.\text{actual_type} == \langle \text{expr} \rangle.\text{expected_type}$
 3. Syntax production: $\langle \text{var} \rangle \rightarrow A \mid B \mid C$
 - Semantic rule: $\langle \text{var} \rangle.\text{actual_type} \leftarrow \text{lookup}(\langle \text{var} \rangle.\text{string})$



Computing Attribute Values

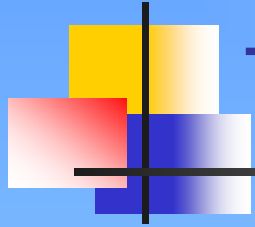
- Let $X_0 \rightarrow X_1 \dots X_n$ be a production
 - If the computing rule of X_0 's attribute is of the form $A(X_0) = f(A(X_1), \dots, A(X_n))$
 - Synthesized attribute
 - If the computing rule of X_j 's attribute is of the form $A(X_j) = f(A(X_0), \dots, A(X_i), \dots, A(X_{j-1}))$, for $i \leq j \leq n$
 - Inherited attribute
- Intrinsic attributes are synthesized attributes of leaf nodes whose values are determined outside the parse tree



A Notation Translation Example Using Syntax-Directed Definition

- $\langle \text{expr} \rangle \rightarrow \langle \text{expr} \rangle + \langle \text{term} \rangle \mid \langle \text{expr} \rangle - \langle \text{term} \rangle \mid \langle \text{term} \rangle$
- $\langle \text{term} \rangle \rightarrow 0 \mid 1 \mid 2 \mid \dots \mid 9$

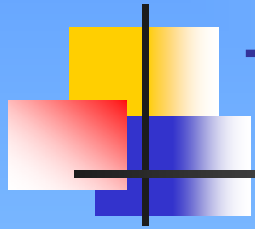
PRODUCTION	SEMANTIC RULES
$\text{expr} \rightarrow \text{expr}_1 + \text{term}$	$\text{expr.t} = \text{expr}_1.t \parallel \text{term.t} \parallel '+'$
$\text{expr} \rightarrow \text{expr}_1 - \text{term}$	$\text{expr.t} = \text{expr}_1.t \parallel \text{term.t} \parallel '-'$
$\text{expr} \rightarrow \text{term}$	$\text{expr.t} = \text{term.t}$
$\text{term} \rightarrow 0$	$\text{term.t} = '0'$
$\text{term} \rightarrow 1$	$\text{term.t} = '1'$
\dots	\dots
$\text{term} \rightarrow 9$	$\text{term.t} = '9'$



Tree Traversals

- Perform depth-first traversal of the parse tree to generate a fully attributed parse tree

```
procedure visit(node N) {  
    for (each child C of N, from left to right) {  
        visit(C);  
    }  
}
```



Translation Schemes

- We used semantic rules as a translation scheme
- Now we use **semantic actions** as a translation scheme to get the same translation result
- Syntax-directed definition for a BNF grammar
 - Associate each grammar symbol (terminals and non-terminals) with a set of attribute
 - Type information for type checking/conversion
 - Notation representation for notation translation
 - Attach a **semantic** rule or program fragment to each production in a grammar
 - Computing the values of the attributes associated with the symbols in the production



New BNF Productions and Parse Trees Using Semantic Actions

- Actions are added in the productions

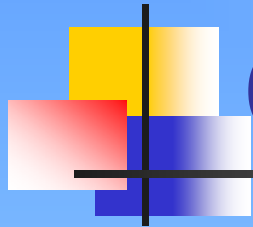
<i>expr</i>	\rightarrow	<i>expr</i> ₁ + <i>term</i>	{print('+')}
<i>expr</i>	\rightarrow	<i>expr</i> ₁ - <i>term</i>	{print('-')}
<i>expr</i>	\rightarrow	<i>term</i>	
<i>term</i>	\rightarrow	0	{print('0')}
<i>term</i>	\rightarrow	1	{print('1')}
		...	
<i>term</i>	\rightarrow	9	{print('9')}

- When drawing a parse tree
 - Indicate an action by constructing an extra child for it, connected by a dashed line to the node that corresponds to the head of the production



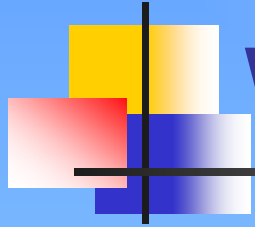
Actions Translating $9-5+2$ into $95-2+$

- Perform a postorder traversal of the parse tree



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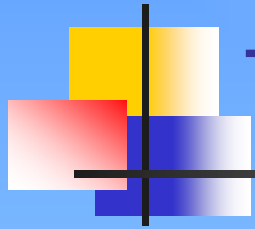
What Can Be Done Now?

- Define language **syntax** using BNF grammar
- Parsing to detect **syntax** errors
 - Syntax analysis
- **Syntax**-directed translation



Define A Language and Syntax-Directed Translation

- $\text{expr} \rightarrow \text{expr} + \text{term} \mid \text{expr} - \text{term} \mid \text{term}$
- $\text{term} \rightarrow 0 \mid 1 \mid \dots \mid 9$
- **Syntax-directed translation based on semantic actions**
- $\text{expr} \rightarrow$
 - $\text{expr} + \text{term} \{ \text{print}('+') \}$
 - $\mid \text{expr} - \text{term} \{ \text{print}('-') \}$
 - $\mid \text{term}$
- $\text{term} \rightarrow$
 - $0 \quad \{ \text{print}('0') \}$
 - $\mid 1 \quad \{ \text{print}('1') \}$
 - $\mid \dots$
 - $\mid 9 \quad \{ \text{print}('9') \}$



Top-Down Parsing

$$A \rightarrow A \alpha \mid \beta$$

$$A \rightarrow \beta R$$

$$R \rightarrow \alpha R \mid \varepsilon$$

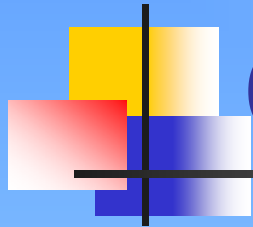
- Left recursion removal for top-down parsing
- $\text{expr} \rightarrow \text{term rest}$
- $\text{rest} \rightarrow + \text{term} \{ \text{print}(' + ') \} \text{rest}$
 $\mid - \text{term} \{ \text{print}(' - ') \} \text{rest}$
 $\mid \varepsilon$
- $\text{term} \rightarrow 0 \{ \text{print}('0') \}$
 $\mid 1 \{ \text{print}('1') \}$
 $\mid \dots$
 $\mid 9 \{ \text{print}('9') \}$



Implementing Parsing and Translation

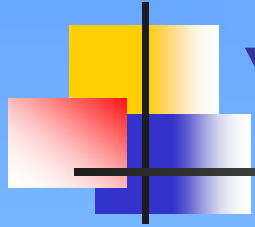
- $\text{expr} \rightarrow$
term rest
- $\text{rest} \rightarrow$
 - + term { print('+') } rest
 - | - term { print('-') } rest
 - | ϵ
- $\text{term} \rightarrow$
 - 0 { print('0') }
 - | 1 { print('1') }
 - | ...
 - | 9 { print('9') }

```
void expr() {  
    term(); rest();  
}  
  
void rest() {  
    if ( lookahead == '+' ) {  
        match('+'); term(); print('+'); rest();  
    }  
    else if ( lookahead == '-' ) {  
        match('-'); term(); print('-'); rest();  
    }  
    else { } /* do nothing with the input */ ;  
}  
  
void term() {  
    if ( lookahead is a digit ) {  
        t = lookahead; match(lookahead); print(t);  
    }  
    else report("syntax error");  
}
```



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You should now be able to ...

- Remove left recursions in BNF;
- Describe syntax-directed definition and attribute grammar;
- Implement a simple language.

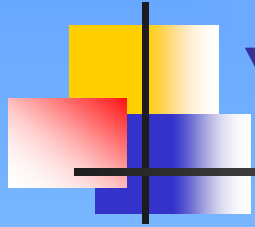


Remove Left Recursion

$$A \rightarrow A \alpha \mid A \beta \mid \gamma$$

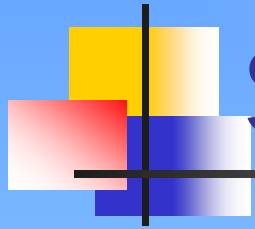
$$A \rightarrow \gamma R$$

$$R \rightarrow \alpha R \mid \beta R \mid \varepsilon$$



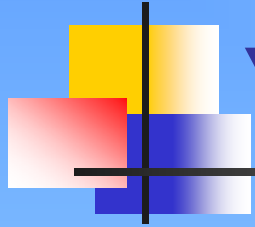
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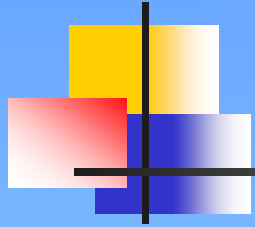
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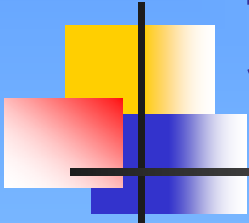
You should now be able to ...

- Remove left recursions in BNF;
- Describe syntax-directed definition and attribute grammar;
- Implement a simple language.



Implement A Simple Language

- Define language **syntax** using BNF grammar
- Parse sentences and detect **syntax** errors
- Use **syntax**-directed definition to perform language translation



Homework (Due on 01/29 at 11:55 PM)

■ 2.1. (20 points)

- (a) Define a BNF grammar for a language that could express a one-digit number, additions and/or subtractions of multiple one-digit numbers in a prefix notation (e.g., $-xy$ is the prefix notation for $x-y$ and the prefix notation of an infix notation $4+5-2+6$ is $+ - + 4 5 2 6$); (5 pts)
- (b) Construct a syntax-directed translation scheme that translates the above-defined one-digit arithmetic expressions from prefix notation into infix notation; (5 pts)
- (c) Implement an executable and correct program to perform the above-mentioned translation. (10 pts)