## TABLE 6 Logical Equivalences.

| Equivalence | Name |
| :---: | :---: |
| $\begin{aligned} & p \wedge \mathbf{T} \equiv p \\ & p \vee \mathbf{F} \equiv p \end{aligned}$ | Identity laws |
| $\begin{aligned} & p \vee \mathbf{T} \equiv \mathbf{T} \\ & p \wedge \mathbf{F} \equiv \mathbf{F} \end{aligned}$ | Domination laws |
| $\begin{aligned} & p \vee p \equiv p \\ & p \wedge p \equiv p \end{aligned}$ | Idempotent laws |
| $\neg(\neg p) \equiv p$ | Double negation law |
| $\begin{aligned} & p \vee q \equiv q \vee p \\ & p \wedge q \equiv q \wedge p \end{aligned}$ | Commutative laws |
| $\begin{aligned} & (p \vee q) \vee r \equiv p \vee(q \vee r) \\ & (p \wedge q) \wedge r \equiv p \wedge(q \wedge r) \end{aligned}$ | Associative laws |
| $\begin{aligned} & p \vee(q \wedge r) \equiv(p \vee q) \wedge(p \vee r) \\ & p \wedge(q \vee r) \equiv(p \wedge q) \vee(p \wedge r) \end{aligned}$ | Distributive laws |
| $\begin{aligned} & \neg(p \wedge q) \equiv \neg p \vee \neg q \\ & \neg(p \vee q) \equiv \neg p \wedge \neg q \end{aligned}$ | De Morgan's laws |
| $\begin{aligned} & p \vee(p \wedge q) \equiv p \\ & p \wedge(p \vee q) \equiv p \end{aligned}$ | Absorption laws |
| $\begin{aligned} & p \vee \neg p \equiv \mathbf{T} \\ & p \wedge \neg p \equiv \mathbf{F} \end{aligned}$ | Negation laws |

TABLE 8 Logical Equivalences Involving Biconditional Statements.

$$
\begin{aligned}
& p \leftrightarrow q \equiv(p \rightarrow q) \wedge(q \rightarrow p) \\
& p \leftrightarrow q \equiv \neg p \leftrightarrow \neg q \\
& p \leftrightarrow q \equiv(p \wedge q) \vee(\neg p \wedge \neg q) \\
& \neg(p \leftrightarrow q) \equiv p \leftrightarrow \neg q
\end{aligned}
$$

TABLE 7 Logical Equivalences Involving Conditional Statements.

$$
\begin{aligned}
& p \rightarrow q \equiv \neg p \vee q \\
& p \rightarrow q \equiv \neg q \rightarrow \neg p \\
& p \vee q \equiv \neg p \rightarrow q \\
& p \wedge q \equiv \neg(p \rightarrow \neg q) \\
& \neg(p \rightarrow q) \equiv p \wedge \neg q \\
& (p \rightarrow q) \wedge(p \rightarrow r) \equiv p \rightarrow(q \wedge r) \\
& (p \rightarrow r) \wedge(q \rightarrow r) \equiv(p \vee q) \rightarrow r \\
& (p \rightarrow q) \vee(p \rightarrow r) \equiv p \rightarrow(q \vee r) \\
& (p \rightarrow r) \vee(q \rightarrow r) \equiv(p \wedge q) \rightarrow r
\end{aligned}
$$

TABLE 1 Rules of Inference.

| Rule of Inference | Tautology | Name |
| :---: | :---: | :---: |
| $\therefore \begin{gathered} p \\ \frac{p \rightarrow q}{q} \end{gathered}$ | $(p \wedge(p \rightarrow q)) \rightarrow q$ | Modus ponens |
| $\begin{aligned} & \neg q \\ & \frac{p \rightarrow q}{\neg p} \end{aligned}$ | $(\neg q \wedge(p \rightarrow q)) \rightarrow \neg p$ | Modus tollens |
| $\begin{aligned} & \quad p \rightarrow q \\ &=\frac{q \rightarrow r}{p \rightarrow r} \end{aligned}$ | $((p \rightarrow q) \wedge(q \rightarrow r)) \rightarrow(p \rightarrow r)$ | Hypothetical syllogism |
| $\begin{aligned} & p \vee q \\ \therefore & \neg p \\ \therefore & q \end{aligned}$ | $((p \vee q) \wedge \neg p) \rightarrow q$ | Disjunctive syllogism |
| $\therefore \frac{p}{p \vee q}$ | $p \rightarrow(p \vee q)$ | Addition |
| $\therefore \frac{p \wedge q}{p}$ | $(p \wedge q) \rightarrow p$ | Simplification |
| $\begin{gathered} p \\ \therefore \frac{q}{p \wedge q} \end{gathered}$ | $((p) \wedge(q)) \rightarrow(p \wedge q)$ | Conjunction |
| $\begin{aligned} & p \vee q \\ & \therefore \neg p \vee r \\ & q \vee r \end{aligned}$ | $((p \vee q) \wedge(\neg p \vee r)) \rightarrow(q \vee r)$ | Resolution |


| TABLE 2 Some Useful Summation Formulae. |  |
| :--- | :--- |
| Sum |  |
| $\sum_{k=0}^{n} a r^{k}(r \neq 0)$ | $\frac{\text { Closed Form }}{r-1}, r \neq 1$ |
| $\sum_{k=1}^{n} k$ |  |
| $\sum_{k=1}^{n} k^{2}$ | $\frac{n(n+1)}{2}$ |
| $\sum_{k=1}^{n} k^{3}$ |  |
| $\sum_{k=0}^{\infty} x^{k},\|x\|<1$ |  |
| $\sum_{k=1}^{\infty} k x^{k-1},\|x\|<1$ | $\frac{n(n+1)(2 n+1)}{6}$ |

## TABLE 1 Set Identities.

| Identity | Name |
| :--- | :--- |
| $A \cap U=A$ | Identity laws |
| $A \cup \emptyset=A$ |  |
| $A \cup U=U$ | Domination laws |
| $A \cap \emptyset=\emptyset$ | Idempotent laws |
| $A \cup A=A$ | Complementation law |
| $A \cap A=A$ | Commutative laws |
| $\overline{(\bar{A})}=A$ |  |
| $A \cup B=B \cup A$ | Associative laws |
| $A \cap B=B \cap A$ | Distributive laws |
| $A \cup(B \cup C)=(A \cup B) \cup C$ |  |
| $A \cap(B \cap C)=(A \cap B) \cap C$ | De Morgan's laws |
| $A \cup(B \cap C)=(A \cup B) \cap(A \cup C)$ |  |
| $A \cap(B \cup C)=(A \cap B) \cup(A \cap C)$ | Absorption laws |
| $\overline{A \cap B}=\bar{A} \cup \bar{B}$ |  |
| $\overline{A \cup B}=\bar{A} \cap \bar{B}$ |  |
| $A \cup(A \cap B)=A$ |  |
| $A \cap(A \cup B)=A$ |  |
| $A \cup \bar{A}=U$ |  |
| $A \cap \bar{A}=\emptyset$ |  |

